# An Energy and Performance Balancing Algorithm Based on Adaptive Cloud Resource Re-Configurability

**[1]Dr. Danthuluri Sudha[1]**
Dept. of. CSE, CMR Institute of Technology, Bengaluru.
**[2]Dr. S Geetha[2]**
[2] Dept. of. CSE, BNM Institute of Technology,  Bengaluru
**[3]Dr N K Vijay Kumar**
, [3]Government Polytechnic,  Mulbagilu..

**Abstract**— Due to its speedier execution and adaptable processing, cloud computing VMs with Dynamic Voltage and Frequency Scaling (DVFS) have achieved tremendous demand in real time. It can meet the major goals of high energy efficiency, resource allocation, and cloud computing VM balancing. In order to properly exploit cloud computing systems, it is important to achieve a high balance between resource allocation and reconfiguration cost in the cloud computing environment. It is vital in cloud computing systems to strike a balance between resource allocation, energy consumption, and information processing center performance, as well as decreasing job load. As a result, an Adaptive Cloud Resource Re- Configurability (ACRR) model is presented here, which relies on Dynamic Voltage and Frequency Scaling (DVFS) to reduce high reconfiguration costs in various data processing applications. Experimental results verify superiority of our proposed ACRR technique in terms of power consumption, average power and power sum.
**Keywords**: DVFS, Cloud Computing, Energy consumption,Performance Balancing.

## I INTRODUCTION

Various companies have changed their attention to applications to the cloud in order to turn a
 profit and effectively allocate resources across all Cloud Computing operations. Cloud Computing is a new approach that promotes the use of resources, data, organisation, storage, and information. To reduce administrative costs, the cloud-computing model connects locations of various editing tools across the cloud organization. Cloud Computing is a newly redesigned organization that offers a higher level of storage limit, faster flexibility, and works on the principle of payment usage, meaning users only pay for the time they spend. The cloud infrastructure applications are classified into infrastructure-as-a- service (IaaS), platform-as-a-service (PaaS), and other cloud platform applications. Virtualization is the most fundamental approach for cloud computing systems to reduce resource usage [1][2]. Virtualization frequently assists by enabling resources on distinct virtual machines (VMs) on a single computer by assigning each asset to a location with the necessary equipment. It is additionally seen that for the most part plenty of resources are allocated to VMs than typically needed [3]. Furthermore, cloud computing systems face very significant task-stacks, which cause

employees to be unable to continue for a lengthy period of time. Demanding are source re-administration course for a set period of time to keep the current resource sharing up to date. [4], [5] [6].

The most important aspect of cloud computing applications is the efficient distribution of resources to extend the exhibition of cloud data centers. In this way, different researchers present various procedures to upgrade the exhibition and address these issues promptly, such as the Hierarchical Reliability-Driven Scheduling (HRDS) method, Constrained Earliest Finish Time (CEFT) algorithm, Contention-Aware Energy-proficient Duplication (FastCEED) strategy, Dynamic Voltage and Frequency Scaling (DVFS), and Voltage and Frequency Scaling (VFS). Dynamic Voltage and Frequency Scaling (DVFS) is perhaps the most widely used technique for scheduling assignment loads. DVFS is a well-established energy consumption advancement model for embedded cloud frameworks, and energy savings can be achieved by dramatically lowering the voltage of any chip. By reducing energy consumption in cloud devices, the DVFS strategy aids in achieving high QoS internet administrations.

## II. RELATED WORK

To improve services and improve the efficiency in heterogenic computing, an effective planning techniques are important, based on the DVFS process, is needed, which provides a variety of power supply and utility offices in achieving better resource utilization and energy efficiency. Mobile-Edge Computing (MEC) is tied in with offering application designers and specialist co-op's cloud-computing abilities and an IT administration at the edge of the versatile organization [7] [8] [9]. Regardless of the way that cloud computing may be used to handle traditional issues, for example, adaptability issues and rapid resource provisioning times, a multifaceted examination is essential when it comes to multi- administrator situations with time-based applications and administrations [10] [11] [12]. Cloud-based mixed media administrations have been broadly utilized. As the developing scale, clients frequently have very different nature of administration (QoS) assumptions. A vital test for separated administrations is the way to ideally dispense cloud resources to fulfill various clients [13].

Cloud computing and virtualization technologies play important roles in today's assistance-oriented computing world. More ordinary administrations are being moved to virtualize computing conditions to accomplish adaptable sending and high accessibility[14] [15]. Presented a timetable algorithm dependent on Fluffy Induction Framework (FIS), for worldwide holder resource allocation by assessing hubs' statuses utilizing FIS. There has been a progressive effort to reduce the energy consumption of large cloud farms through increased utilization of capacity and load- fixing methods. With the introduction of the new Container as a service (CaaS) by cloud providers, maximizing the use of virtual machine (VM) technology is very important [16] [17]. We have broken the cloud tracking logs from the Google team and looked at the remaining dynamic cloud functionality, which is critical to testing and approving our proposed systems. The test results showed a 7.55% improvement in standard power consumption compared to standard conditions where the size of the machine size was

set. Similarly, compared to normal conditions, the total number of VMs launched to assist management is further enhanced by 68% overall [18] [19] [20].

Arranging multiprocessors has two main objectives: energy efficiency and improved efficiency [21] [22]. This study proposes Contention-aware, Energy-Efficient, Duplication- based-Mixed Integer Programming (CEEDMIP) details of the activity diagram that organizes various multiprocessors linked to a distributed system or organization in chip construction. Compared to other duplicate power algorithms, the developed MIP with a bunching-based heuristic gives flexibility and an increase in strength by 10-30% with improved time and accuracy. The issues of energy-obliged and time-compelled task scheduling on numerous heterogeneous PCs are examined as combinatorial improvement issues [23] [24]. We can locate an ideal parcel of a given remaining burden and utilize adjusted rundown scheduling (MLS) algorithm to produce a segment of the arrangement of errands that is an estimation of the ideal outstanding task at hand segment.

Reducing energy utilization has become a significant objective in planning current heterogeneous computing frameworks. [25] This paper tends to execution issues of resource allocation in cloud computing. Following this perception, we propose another correspondence model of cloud computing application, called CA-DAG. This model depends on Directed Acyclic Graphs that not with standing computing vertices incorporate separate vertices to address correspondences. The proposed CA-DAG model makes space for the advancement of various existing answers for resource allocation and for creating novel scheduling plans of improved productivity [26 To address these concerns, we created a Dynamic Voltage and Frequency Scaling (DVFS) mechanism for cloud computing devices based on Adaptive Cloud Resource Re-Configurability (ACRR) that efficiently minimizes power usage, while performing tasks in a very short time. In heterogeneous computing gadgets, our proposed ACRR approach aids in achieving a high compromise between energy and execution. By focusing on lower transmission rates and minimizing deferral, this scheduling strategy aids in achieving high QoS. This strategy effectively aids in achieving the goal of lowering the cost of communication and computing resources. This approach adapts to changing situations through organizing programs, allocating resources, and allowing for flexible implementation. In comparison to other existing methods, our proposed ACRR method performs significantly better.

## II .PROPOSED ENERGY BALANCED SCHEDULING ARCHITECTURE

The proposed ACRR method and its different modules are described as follows. This segment also discusses the reduction of computational and re-design costs in datacenters. The proposed ACRR approach is depicted in Figure 1.1. For cloud computing VMs, we present a novel Adaptive Cloud Resource Re-Configurability (ACRR) approach. The proposed ACRR approach is based on the concept that it can accommodate a broad variety of cloud computing VMs limited by a central resource controller. Each cloud computing VM, as a self-supervising processor, completes the currently distributed task without the need for anybody else to manage its storage and assets. A distributed programming approach is employed for communication. When a

new task is issued, a central resource controller begins deploying resources that effect the modules while also providing confirmation oversight.

As shown in Figure 1.1, our proposed ACRR procedure has three critical segments that assist with achieving better resource utilization, like data storage, exchanged Local Area Network (LAN), and Virtual Machine Handler (VMC). When some work is assigned to the virtual machine, the size of the task is indicated in smaller segments. The absolute handling season of relegated task is not exactly or equivalent to the assessed utilized time which is extremely fundamental for any strategy to be received progressively situations. Our suggested ACRR technique includes the errand granularity, which represents the maximum amounts of work that could be integrated into the reduced work. Accepting the greatest VMs in relegated assignments using our proposed ACRR strategies can be conveyed by $N_J \square 1$ and illustrated in Figure 1.
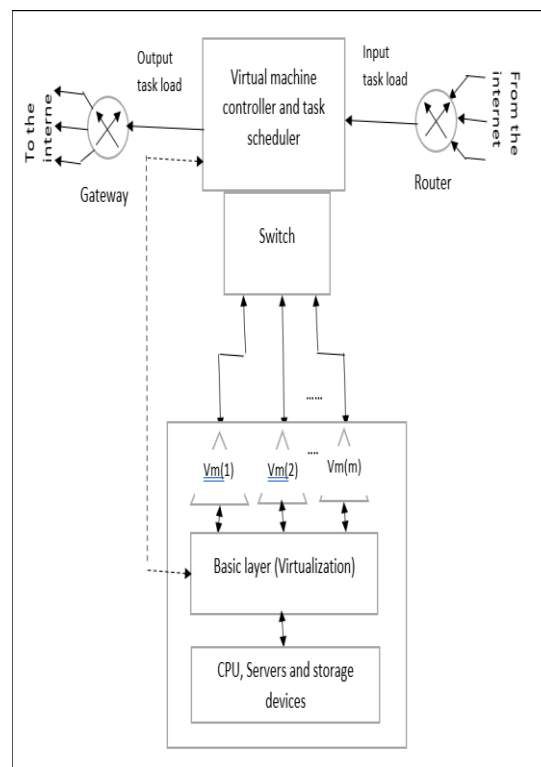


**Figure 1. Proposed ACRR model's architecture diagram**

The rule that each VM is capable of measure $i_g$ $bits$ is being used at by our scheduling method (Bits per Second). Depending on the size S, the working rate $i_g$ will be scaled in sections in parallel at the time of execution. Accept that all errands follow the stretch $[0,{}^t_g]$, with $i^t_g$ occupying the position with the highest working rate. Furthermore, the task size S does not affect the approximate time it will take VM to complete the assigned mission, which was previously set to our model in continuous situations and denoted in a second or two. A VM can also manage the work-heaps and task-load size $S_{\mathbb{b}}$ of a currently assigned task of size S. The OS (Operating System) programs are responsible for the foundation task-loads. The foundation task-load is supposed to be stored in the VM's fundamental memory. As a result, the foundation task-load only necessitated computational costs and did not necessitate communication costs. The consumption

boundary $\beta$ can be communicated in this way.

$$\beta \triangleq i_g .(i^t_g)^{-1} \in [0,1], \qquad (1)$$

The dynamic aspects of computing resources are the most important factor in lowering the computational cost, as shown in equation (1). Assume that the total energy needed by VM to complete a single task with a time interval of ($\mathbb{d}$) is denoted by $\varphi_g$ in joule at the operating rate $i_g$. As a consequence, the dimensionless ratio can be written as

$$\theta(\beta) \triangleq \varphi(i)_g.(\varphi^t_g)^{-1} \equiv \theta(i)_g.(i^t_g)^{-1}, \qquad (2)$$

Where equation (2) represents the total energy consumption of the concerned VM. The DVFS-based CPU analytical method, for example, can be expressed using the equation below.

$$(\beta)=\beta^2, \qquad \beta \epsilon [0,1], \qquad (3)$$

We may also use $\mu$ to measure the relative energy cost of at ask performed by the concerned VM.

## A. Task-Load Reduction using Proposed *ACRR* Technique

In this section, we'll talk about task-load reduction modelling. Assume that $N \triangleq \downarrow$ $\{N_J\}$ is the number of non-overlapping assignments that can be completed in parallel. Assume $S_d$ is the task size given by the computing VM (d). The processing time of different tasks is unaffected by the task length $S_d$. As a result, the speed of processing can be measured in bits per second.

$$i_g(d)=S_d.(\omega)^{-1}, \qquad (4)$$

Equation (4) indicates the maximum duration $S_d^\uparrow = \omega.i_g^\uparrow(d)$. And $\beta \triangleq i_g(d).\left(i_g^\uparrow\right)^{-1} \equiv S_d.\left(S_d^\uparrow\right)^{-1}$. The total size of a job is $S_a$ in bits, and the task size is $S_d \gg 0, d = 1, \dots, N$, which is assigned to the $VM(d)$ by the task scheduler as shown in figure 1.1. We divide the total job size $S_a$ into $N$ parallel tasks to reduce work loads, with $\sum_d^N S_d = S_a$ an as the size boundary limit.

## B. Optimization of Reconfiguration Cost Using Proposed ACRR Technique:

The streamlining of reconfiguration cost is clearly shown in this section. The VM module regulator is used to change the task burdens and monitor Virtual Machines, among other things. The Virtual Machine Controller (VMC), which assists with the last planning of VM resources on various computing VMs, is necessary for controlling the Virtualization Layer as shown in figure 1. The condition (9) can be used to represent the

VM's quality boundaries,

$$\{\omega, i_g^\uparrow(d), \theta_d(\beta_d), \varphi_g^\uparrow(d), \mu(d), S_\mathbb{b}(d), \qquad d = 1, \dots \dots N \ \} \qquad (5)$$

As shown in Figure 1.1, the virtualization layer will represent each of these boundaries before sending them to the Virtual Machine Controller (VMC). The working rate $i_g$ can be scaled up or downsized utilizing a proficient recurrence scaling plan which is constrained by VMC. The force utilization while changing from working recurrence $i_1$ to frequency $i_2$ can be $\varphi(i_1 : i_2)$ in joule. This force utilization basically depends up on the strategy utilized in addition to the workstation's CPUs. This function $\varphi(i_1 : i_2)$ comprises of certain properties, for example, the function $\varphi(i_1 : i_2)\varphi$ depend upon the whole frequency gap$|i_1 - i_2|$ , it gets zero at $i_1 = i_2$ and stay non-diminishing in the whole frequency gap $|i_1 - i_2|$, it is consolidated in a curved manner at $i_1$ and$i_2$. Our model ACRR have a few attributes which can be shown utilizing condition (10),

$$\varphi(i_1 : i_2) = \mathbb{e}_f (i_1 - i_2)^2, \qquad joule \qquad (6)$$

Where,$\mathbb{e}_f$ addresses the cost of reconfiguration unit exchanging of recurrence and the estimations of $\mathbb{e}_f$ is limited distinctly to approximately many $\mu joules/(MHz)^2$. In our model ACRR, for each work the size $S_a$ stays same throughout the individual working time $E_a$ and any sort of changes not happened in the undertaking during execution, there are a lot of loads.Since the prompted time overhead for DVFS-empowered models using the recurrence scaling strategy is extremely low (a few seconds), different tasks can be equally performed at run-time.

The previously stated prediction that the usage boundary $\beta$ can be esteemed indefinitely and that it needs constant computational speeds, as indicated by $i_g$ , is right. Virtual machines (VMs) may have a limited collection of CPUs.

$$\mathbb{?} \triangleq \{\hat{\imath}^{(0)} \equiv 0, \hat{\imath}^{(1)}, \dots, \hat{\imath}^{(\mathbb{L}-1)} \equiv i_g^\uparrow\}, \qquad (7)$$

The discrete computing rates $\mathbb{L}$ make up these finite sets. Both continuous and discrete DVFS enabled methods will benefit from the use of equation (8) to eliminate the optimality loss.

$$\mathbb{F} \triangleq \{\hat{\beta}^{(0)} \equiv 0, \hat{\beta}^{(1)}, \dots \dots, \hat{\beta}^{(\mathbb{L}-1)} \equiv 1\}, \qquad (8)$$

In this case, the specific value set of $\beta$ that denotes the frequency range $\mathbb{?}$ as shown in equation (7). A virtual power consumption curve could be written in $\tilde{\theta}(\beta)$ and is composed of *piecewise linear interpolation* as well as the licensed operating points,

$$\left\{\left(\hat{\beta}^{(\mathbb{o})}, \theta\left(\hat{\beta}^{(\mathbb{o})}\right)\right), \mathbb{o} \right. \qquad (9)$$
$$= 0, \dots \dots, (\mathbb{L})$$
$$\left. - 1\right\}$$

In this case, the relevant vertex points can be seen as,

$$\left(\hat{\beta}^{(\omega)}, \theta\left(\hat{\beta}^{(\omega)}\right)\right) \ and \ \left(\hat{\beta}^{(\omega+1)}, \theta\left(\hat{\beta}^{(\omega+1)}\right)\right) \tag{10}$$

These are previously stated. virtual power consumption curve keeps up the congruity that will be utilized and providing virtual machine services. Similarly, utilization in piecewise linear interpolation proposes that with the assistance of virtual force utilization bend, the normal DVFS's electricity cost empowered procedures stay below the assessed time period term. Here, each VM setup depend based on the CPU type, memory capacity and time cost span. That VMs price depend on the kind for setup.

## IV. Performance Evaluation

Because of the widespread use of enlightening gadgets, advanced instruments, network machines, and compact devices, the demand for cloud computing gadgets has steadily increased. The mixed media signal-preparing strategy is a notable method that can be used in these cloud computing gadgets. As a result of the widespread interest in computing devices in everyday life, the presentation of these devices should be improved. Nonetheless, these computing devices' high energy consumption can degrade their performance. As a result, this section looks at the balancing of execution and energy utilization. To accomplish these targets, we have presented a Dynamic Voltage and Frequency Scaling (DVFS) based Adaptive Cloud Resource Re-Configurability (ACRR) strategy for heterogeneous computing gadgets which proficiently lessens energy utilization just as give predominant execution. The run-time can be computed for different positions such as 25, 50, 100, and 1000. In addition, we present a graphical representation of our findings, which takes into account execution time, task number, and energy consumption. The different limits in Table 1.1, which is shown in the following section, can be used to determine the amount of run-time and total energy consumed. On the Montage scientific dataset, we tested our proposed ACRR model. We found rational workflow tests of different sizes, including 25, 50, 100, and 1000. Our proposed ACRR model runs on a 64-cycle Windows 10 operating system with 16 GB of RAM and an INTEL (R) center (TM) i5-4460 processor. It has a 3.20 GHz processor. Eclipse WS Neon is used to recreate this project.

## a. Comparative Study

In this cutting-edge period, computing gadgets have administered market in various fields like clinical, medical care arrangements, exchanging, programming organizations and enterprises, and so forth. In this way, future aptitude is plainly for these cloud computing gadgets because of their broad necessities. The effectiveness of these computing gadgets might be diminished because of high energy utilization and the absence of proficient resource use strategies. Therefore, these issues can be figured out utilizing effective assignment scheduling methods. Accordingly, to assign resources appropriately and plan all the undertakings proficiently to conquer power utilization issues, we have introduced a novel Dynamic Voltage and Frequency Scaling (DVFS)

based Adaptive Cloud Resource Re-Configurability (ACRR) strategy. Table 1.1 compares the results of using the logical model montage for various positions such as 25, 50, 100, and 1000 in comparison to other state-of-the-art methods in terms of energizing the framework, increasing relationships with endorsers, offering better resource use, and bearing to deal with multiple errands all at once, and so on. Table 1.1 illustrates the average power, power consumption, and power sum for the logical model montage using our ACRR technique, as well as comparisons to other state-of-the-art methods at various positions such as 25, 50, 100, and 1000.

**b. Graphical Representation**

This section displays a graphical representation of the results of our analysis. Figure 1.2 shows how our proposed ACRR strategy compares to the DVFS approach for different positions such as 25, 50, 100, and 1000, using montage as the logical outstanding task. For positions 25, 50, 100, and 1000, Figure 1.2 provides a Power Sum Comparison of our proposed ACRR method with a DVFS approach focused on logical remaining burden montage. For various positions such as 25, 50, 100, and 1000, Figure 1.3 compares the Average Power required of our proposed ACRR method with the DVFS procedure using logical outstanding burden montage. For different positions such as 25, 50, 100, and 1000, Figure 1.4 compares the power consumption of our proposed ACRR system with the DVFS, Interquartile Range (IQR), Median Absolute Deviation (MAD), Local Regression (LR) process using logical outstanding burden montage. This result shows the prevalence of our proposed ACRR strategy in terms of power sum, average power, and power consumption.

**TABLE 1.1 Comparison of various parameters for ACRR Technique using Scientific Model Montage**

| Algorithm | Montage | Parameters | | |
|---|---|---|---|---|
| | | Power Sum (W) | Average Power (W) | Power Consumption (Wh) |
| IQR | 25 | 537802.69 | 27.12 | 173.10 |
| | 50 | 1193231.28 | 27.12 | 469.70 |
| | 100 | 2516667.71 | 27.12 | 1294.47 |
| | 1000 | 26453383.01 | 27.12 | 73998.96 |
| MAD | 25 | 526941.52 | 25.08 | 169.12 |
| | 50 | 1169133.45 | 25.08 | 458.89 |
| | 100 | 2465842.50 | 25.08 | 1255.15 |
| | 1000 | 25919145.31 | 25.08 | 71594.76 |
| LR | 25 | 521235.66 | 24.95 | 164.80 |
| | 50 | 1156473.76 | 24.95 | 446.18 |
| | 100 | 2439141.70 | 24.95 | 1232.99 |
| | 1000 | 25638485.87 | 24.95 | 70668.49 |

| | | | | |
|---|---|---|---|---|
| DVFS | 25 | 500514.37 | 28.65 | 159.32 |
| | 50 | 1110499.1 | 28.65 | 432.30 |
| | 100 | 2342175.63 | 28.65 | 1191.41 |
| | 1000 | 24619248.96 | 28.65 | 68107.65 |
| ACRR | 25 | 126978.33 | 21.99 | 39.18 |
| | 50 | 157883.18 | 21.99 | 52.60 |
| | 100 | 249275.55 | 21.99 | 95.82 |
| | 1000 | 1922348.05 | 21.99 | 937.90 |
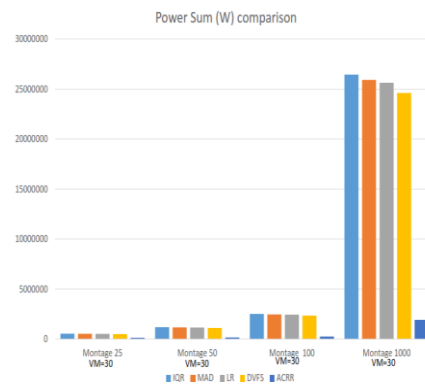


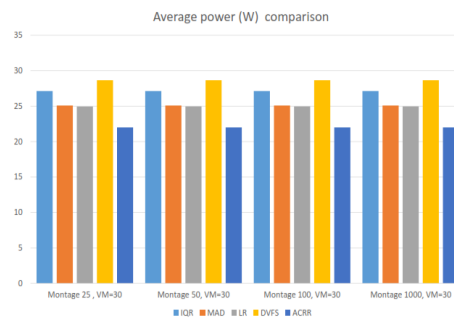**Figure 1.2 Power Sum comparison using theACRR technique**



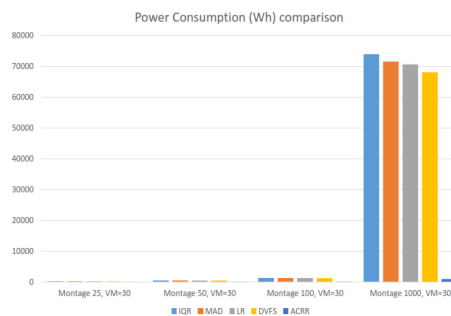**Figure 1.3 Average Power comparisons using theACRR technique**



**Figure 1.4 Power Consumption comparison using our ACRR technique**

**V Conclusion**

For cloud computing VMs, managing power dissipation and maintaining a large utilization among appropriate resource distribution and minimized reconfiguring costs are critical. As a result, an Adaptive Cloud Resource Re-Configurability (ACRR) solution based on Dynamic Voltage and Frequency Scaling (DVFS) is used for cloud computing VMs, which effectively reduces energy consumption while still executing operations in a short amount of time. Because of its high speed and efficient energy reduction, the proposed ACRR technique with IQR, MAD, LR, and DVFS helps to minimize task load and achieve better resource utilization. An efficient modelling approach for reducing reconfiguration costs, reducing task load, and improving performance is presented. The experimental results are compared with IQR, MAD, LR and DVFS in terms of average power, energy consumption, and power sum. Montage scientific dataset for Montage 25 is 39.18879852 Watts, Montage 50 is 52.60213278 Watts, Montage 100 is 95.8242739 Watts, and Montage 1000 is 937.9076791 Watts, which is very less in comparison with IQR, MAD, LR and DVFS and concludes high superiority of the proposed ACRR technique.

**References**

1. Y. Zhang, J. Ren, J. Liu, C. Xu, H. Guo, and Y. Liu, "A survey on emerging computing paradigms for big data," Chinese Journal of Electronics, vol. 26, no. 1, pp. 1–12, 2017.
2. G. Juve, A. Chervenak, E. Deelman, S. Bharathi, G. Mehta, and K. Vahi, "Characterizing and profiling scientific workflows," Future Generation Computer Systems, vol. 29, no. 3, pp. 682–692, 2013.
3. J. Ren, H. Guo, C. Xu, and Y. Zhang, "Serving at the edge: A scalable iot architecture based on transparent computing," IEEE Network, vol. 31, no. 5, pp. 96–105, 2017.
4. L. T. Yang and M. Guo, High-performance computing: paradigm and infrastructure. John Wiley & Sons, 2005, vol. 44.
5. V. K. Naik, C. Liu, L. Yang, and J. Wagner, "Online resource matching for heterogeneous grid environments," in ccgrid. IEEE, 2005, pp. 607– 614.
6. Y.-K. Kwok and I. Ahmad, "Static scheduling algorithms for allocating directed task graphs to multiprocessors," ACM Computing Surveys (CSUR), vol. 31, no. 4, pp. 406–471, 1999.
7. D. P. Vidyarthi, B. K. Sarker, A. K. Tripathi, and L. T. Yang, "Scheduling in distributed computing systems: Analysis, design and models". Springer Science & Business Media, 2008.
8. M. Lin and L. T. Yang, "Hybrid genetic algorithms for scheduling partially ordered tasks in a multi-processor environment," in Real-Time Computing Systems and Applications, 1999. RTCSA'99. Sixth International Conference on. IEEE, 1999, pp. 382–387.
9. S. Abrishami, M. Naghibzadeh, and D. H. Epema, "Deadline-constrained workflow scheduling algorithms for infrastructure as a service clouds," Future Generation Computer Systems, vol. 29, no. 1, pp. 158–169, 2013.

10. Z. Cai, X. Li, R. Ruiz, and Q. Li, "A delay-based dynamic scheduling algorithm for bag-of-task workflows with stochastic task execution times in clouds," Future Generation Computer Systems, vol. 71, pp. 57–72, 2017.

11. D. Zhang, Z. Chen, L. X. Cai, H. Zhou, S. Duan, J. Ren, X. Shen, and Y. Zhang, "Resource allocation for green cloud radio access networks with hybrid energy supplies," IEEE Transactions on Vehicular Technology, vol. 67, no. 2, pp. 1684–1697, 2018.

12. K. Li, X. Tang, B. Veeravalli, and K. Li, "Scheduling precedence constrained stochastic tasks on heterogeneous cluster systems," IEEE Transactions on computers, vol. 64, no. 1, pp. 191–204, 2015.

13. A. Hameed, A. Khoshkbarforoushha, A. Zomaya, "A survey and taxonomy on energy efficient resource allocation techniques for cloud computing systems", Springer-Verlag Wien: June2014.

14. X. Zhu, R. Ge, J. Sun, and C. He, "3e: Energy-efficient elastic scheduling for independent tasks in heterogeneous computing systems," Journal of Systems and Software, vol. 86, no. 2, pp. 302– 314, 2013.

15. D. Kliazovich, J. E. Pecero, A. Tchernykh, P. Bouvry, S. U. Khan, and A. Y. Zomaya, "Ca-dag: Modeling communication-aware applications for scheduling in cloud computing," Journal of Grid Computing, vol. 14, no. 1, pp. 23–39, 2015.

16. Calheiros RN, Ranjan R, Beloglazov A, De Rose CAF, Buyya R. "CloudSim: A Toolkit for Modeling and Simulation of Cloud Computing Environments and Evaluation of Resource Provisioning Algorithms". Softw Pract Exper. 2011; 41-50.

17. Belogazov A, Abawajy J, Buyya R. "Energy-aware resource allocation heuristics for efficient management of data centers for Cloud computing[J]". Future Generation Computer Systems,2012, 28(5):755-768.

18. Gai K, Qiu M, Zhao H, et al. "Dynamic energy-aware cloudlet-based mobile cloud computing model for green computing[J]". Journal of Network & Computer Applications, 2016, 59I:46-54.

19. S. Selvarani, A. Julian and E. I. Nehru, "Enhanced two stage heuristics algorithm for VM scheduling," *2016 International Conference on Advanced Communication Control and Computing Technologies (ICACCCT)*, Ramanathapuram, 2016, pp. 726-729

20. E.Iniya Nehru, Anith Julian, Selvarani.S, "Two stage optimal VM Scheduling scheme for IaaS Cloud" in Proc. International Conference on Communication and Security , 2016, Pondicheerry Engineering College.

21. G. Xie, Y. Chen, Y. Liu, Y. Wei, R. Li and K. Li, "Resource Consumption Cost Minimization of Reliable Parallel Applications on Heterogeneous Embedded Systems," in *IEEE Transactions on Industrial Informatics*, vol. 13, no. 4, pp. 1629-1640, Aug. 2017.

22. Z. Zhu, G. Zhang, M. Li and X. Liu, "Evolutionary Multi-Objective Workflow Scheduling in Cloud," in *IEEE Transactions on Parallel and Distributed Systems*, vol. 27, no. 5, pp. 1344-1357, May 1 2016.

23. Prasanna K.P, Sreekanth A.M, "Data Sharing with Sensitive information hiding in data storage using Cloud Computing", International Journal of Trend in Scientific Research and Development, Volume-4, Issue-2, Feb 2020.

24. Stephen, "Energy Saving by Migrating Virtual Machine to Green Cloud Computing", International Journal of Trend in Scientific Research and Development, Volume-4, Issue-3, April 2020.

25. Bryan Harris and Nihat Altiparmak "Monte Carlo Based Server Consolidation for Energy Efficient Cloud Data Centers", 2019 IEEE International Conference on Cloud Computing Technology and Science (CloudCom), 2330-2186.

26. Lei Shi, Jing Xu, Lunfei Wang, Jie Chen, ZhifengJin, Tao Ouyang, Juan Xu, Yuqi Fan, "Multijob Associated Task Scheduling for Cloud Computing Based on Task Duplication and Insertion", Wireless Communications and Mobile Computing, vol. 2021, Article ID 6631752, 13 pages, 2021.

27. Padmavathi, M., and Shaik Mahaboob Basha. "A Conceptual Analysis on Cloud Computing ACO Load Balancing Techniques." International Journal of Computer Science and Engineering (IJCSE) 6.4 (2017): 39-48.

28. Bhatta, Goutam. "Cloud computing." International Journal of Computer Science and Engineering (IJCSE) 5.3 (2016): 7-16.